

# Leverage Scores: Sensitivity and an App

Ilse Ipsen

Joint work with Thomas Wentworth

North Carolina State University  
Raleigh, NC, USA

Research supported by NSF CISE CCF, NSF DMS, DARPA XData

## Leverage Scores

- Real  $m \times n$  matrix  $\mathcal{A}$  with  $\text{rank}(\mathcal{A}) = n$
- Columns of  $Q$  are orthonormal basis for  $\text{range}(\mathcal{A})$   
column  $i$  of  $Q \perp$  column  $j$  of  $Q$ ,  $\|\text{column } j \text{ of } Q\|_2 = 1$
- Leverage scores of  $\mathcal{A}$

$$\ell_j(\mathcal{A}) = \|\text{row } j \text{ of } Q\|_2^2 \quad 1 \leq j \leq m$$

# Leverage Scores

- Real  $m \times n$  matrix  $\mathcal{A}$  with  $\text{rank}(\mathcal{A}) = n$
- Columns of  $Q$  are orthonormal basis for  $\text{range}(\mathcal{A})$   
column  $i$  of  $Q \perp$  column  $j$  of  $Q$ ,  $\|\text{column } j \text{ of } Q\|_2 = 1$

- Leverage scores of  $\mathcal{A}$

$$\ell_j(\mathcal{A}) = \|\text{row } j \text{ of } Q\|_2^2 \quad 1 \leq j \leq m$$

- $0 \leq \ell_j(\mathcal{A}) \leq 1 \quad \sum_{j=1}^m \ell_j(\mathcal{A}) = n = \|\mathcal{A}\|_F^2$
- Outlier detection in least squares/regression problems  
[Hoaglin & Welsch 1978, Velleman & Welsch 1981, Chatterjee & Hadi 1986]
- Importance sampling in randomized matrix algorithms  
[Avron, Boutsidis, Drineas, Mahoney, Toledo, ...]

# Largest Leverage Score

Largest leverage score = coherence of  $\mathcal{A}$

$$\mu(\mathcal{A}) \equiv \max_j \ell_j(\mathcal{A})$$

[Donoho & Ho 2001, Candés, Romberg & Tao 2006, Candés & Recht 2009, ...]

# Largest Leverage Score

Largest leverage score = coherence of  $\mathcal{A}$

$$\mu(\mathcal{A}) \equiv \max_j \ell_j(\mathcal{A})$$

[Donoho & Ho 2001, Candés, Romberg & Tao 2006, Candés & Recht 2009, ...]

- $n/m \leq \mu(\mathcal{A}) \leq 1$
- **Uniform** leverage scores  $\ell_j(\mathcal{A}) = n/m$  for **all**  $j$   
 $\Rightarrow$  **Minimal** coherence  $\mu(\mathcal{A}) = n/m$   
Sampling is **easy**
- **Large** leverage score  $\ell_j(\mathcal{A}) = 1$  for **some**  $j$   
 $\Rightarrow$  **Maximal** coherence  $\mu(\mathcal{A}) = 1$   
Sampling is **hard**

# Why Leverage Scores?

- Analysis of randomized algorithms:  
Quantify the **difficulty of sampling**
- **Probabilities** in randomized matrix algorithms  
[Boutsidis, Drineas, Mahoney, ...]
- Least Squares Solver *Blendenpik* [Avron, Maymounkov, Toledo 2010]  
Bounds for **condition numbers of sampled matrices**

# Why Leverage Scores?

- Analysis of randomized algorithms:  
Quantify the **difficulty of sampling**
- **Probabilities** in randomized matrix algorithms  
[Boutsidis, Drineas, Mahoney, ...]
- Least Squares Solver *Blendenpik* [Avron, Maymounkov, Toledo 2010]  
Bounds for **condition numbers of sampled matrices**

**Condition number** (sensitivity of basis to perturbations)

If  $\mathcal{A}$  has linearly independent columns then

$$\kappa(\mathcal{A}) \equiv \|\mathcal{A}\|_2 \|\mathcal{A}^\dagger\|_2$$

If  $Q$  has **orthonormal columns** then  $\kappa(Q) = 1$

# Overview

- 1 Motivation
- 2 The App
- 3 Sensitivity of leverage scores, to:  
*Rotation of subspace*  
*Matrix perturbations*



# Motivation

# The Problem

Given

Real  $m \times n$  matrix  $Q$  with orthonormal columns  
coherence  $\mu$  and leverage scores  $\ell_j$

Real  $c \times m$  "sampling" matrix  $S$  with  $n \leq c \ll m$

Condition number of sampled matrix

$$\kappa(SQ) \equiv \|SQ\|_2 \|(SQ)^\dagger\|_2$$

Sensitivity of  $SQ$  to perturbations

Given  $\eta$  and  $\delta$ , for which values of  $c$  (in terms of  $\mu$  and  $\ell_j$ ) is

$$\kappa(SQ) \leq 1 + \eta$$

with probability at least  $1 - \delta$ ?

## Example: A Bound [Ipsen & Wentworth 2012]

Label leverage scores so that  $\mu \equiv \ell_{[1]} \geq \dots \geq \ell_{[m]}$

$$\tau \equiv \sum_{j=1}^t \ell_{[j]} + \left( \frac{1}{\mu} - t \right) \ell_{[t+1]} \quad \text{where } t \equiv \lfloor 1/\mu \rfloor$$

SQ :  $c$  rows of  $Q$  sampled **uniformly with** replacement

## Example: A Bound [Ipsen & Wentworth 2012]

Label leverage scores so that  $\mu \equiv \ell_{[1]} \geq \dots \geq \ell_{[m]}$

$$\tau \equiv \sum_{j=1}^t \ell_{[j]} + \left(\frac{1}{\mu} - t\right) \ell_{[t+1]} \quad \text{where } t \equiv \lfloor 1/\mu \rfloor$$

SQ :  $c$  rows of  $Q$  sampled **uniformly with** replacement

Given  $\epsilon$ , with probability at least  $1 - \delta$

$$\kappa(SQ) \leq \sqrt{\frac{1 + \epsilon}{1 - \epsilon}}$$

provided  $c \geq m \mu (2\tau + \frac{2}{3}\epsilon) \ln(2n/\delta)/\epsilon^2$

## Example: A Bound [Ipsen & Wentworth 2012]

Label leverage scores so that  $\mu \equiv \ell_{[1]} \geq \dots \geq \ell_{[m]}$

$$\tau \equiv \sum_{j=1}^t \ell_{[j]} + \left(\frac{1}{\mu} - t\right) \ell_{[t+1]} \quad \text{where } t \equiv \lfloor 1/\mu \rfloor$$

$SQ$  :  $c$  rows of  $Q$  sampled uniformly with replacement

Given  $\epsilon$ , with probability at least  $1 - \delta$

$$\kappa(SQ) \leq \sqrt{\frac{1 + \epsilon}{1 - \epsilon}}$$

provided  $c \geq m \mu (2\tau + \frac{2}{3}\epsilon) \ln(2n/\delta)/\epsilon^2$

In practice:  $\kappa(SQ) \leq 10$  with at least 99% probability  
provided  $c \geq m \mu (2.1\tau + .7) (\ln(2n) + 4.7)$

## Practical Questions

For various values of  $m$ ,  $n$ , coherence, leverage scores, and  $c$ :

- How does this bound compare to existing bounds?
- How does this bound compare to the condition numbers of sampled matrices?
- What are the smallest values of  $c$  for which the bound becomes informative?
- What are the smallest values of  $c$  for which the sampled matrices have full rank?
- How large are the condition numbers of the sampled matrices?
- When is sampling 10% of the rows sufficient?
- For a given coherence, are there leverage score distributions that make sampling harder?
- How large can  $n/m$  be, before sampling becomes inefficient?

# The App

# Matlab App with GUI (arXiv:1402.0642)

$\kappa_{SQ} = \kappa(SQ)$

- 1 Plot *probabilistic bounds* for  $\kappa(SQ)$
- 2 Run *experiments*, and plot *actual values* of  $\kappa(SQ)$



# Matlab App with GUI (arXiv:1402.0642)

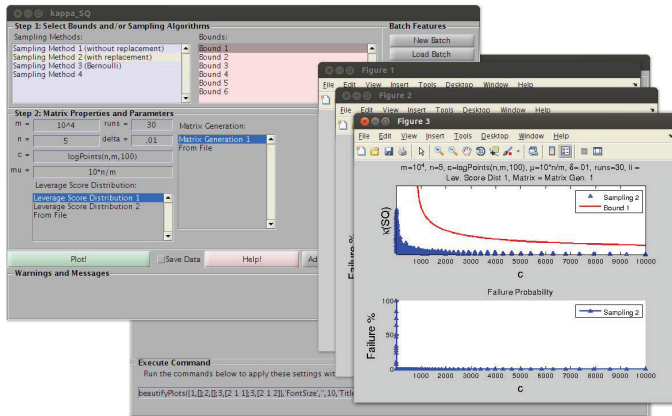
$\text{kappaSQ} = \kappa(SQ)$

- 1 Plot *probabilistic bounds* for  $\kappa(SQ)$
- 2 Run *experiments*, and plot *actual values* of  $\kappa(SQ)$

## Features of kappaSQ

- Four randomized sampling methods
  - Uniform & leverage score sampling with replacement*
  - Uniform sampling without replacement*
  - Bernoulli sampling*
- Six probabilistic bounds
- Test matrix generation (for given  $m$ ,  $n$  and leverage scores)
- Leverage score distributions: "adversarial" or not  
(for given  $m$ ,  $n$  and coherence)
- "Publication-ready" plots
- Easy incorporation of user's own codes

# Screen Shot of kappaSQ



# Advantages of kappaSQ

- Insight into behavior of sampling methods & bounds  
in a practical, non-asymptotic context

*Compare different bounds*

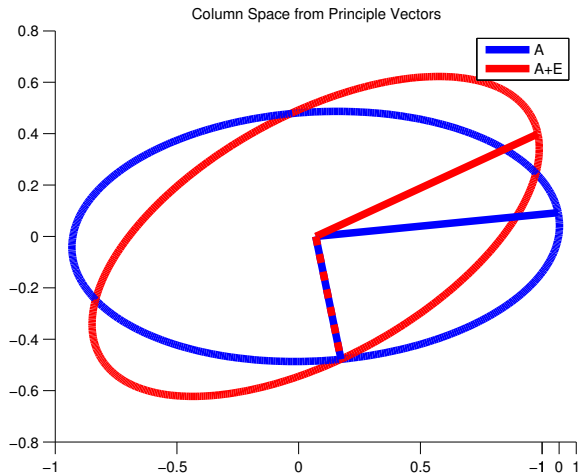
*Compare bounds to experiments*

*Explore the limits of randomized sampling*

- Intuitive user interface (plot button)
- Visually appealing plots
- Sensible default values
- Extensive facilities for customizing plots (beautify)
- Very little familiarity with Matlab required

## Sensitivity of Leverage Scores to Rotation of Subspace

# Rotation of Subspace



## Exact and Perturbed Leverage Scores

- Exact subspace:  $\text{range}(\mathcal{A})$ ,  $\mathcal{A}$  is  $m \times n$  with  $\text{rank}(\mathcal{A}) = n$   
Exact leverage scores

$$\ell_j(\mathcal{A}) = \|e_j^T \mathcal{A}\|_2^2 \quad 1 \leq j \leq m$$

where  $A$  is orthonormal basis for  $\text{range}(\mathcal{A})$

## Exact and Perturbed Leverage Scores

- Exact subspace:  $\text{range}(\mathcal{A})$ ,  $\mathcal{A}$  is  $m \times n$  with  $\text{rank}(\mathcal{A}) = n$   
Exact leverage scores

$$\ell_j(\mathcal{A}) = \|e_j^T \mathcal{A}\|_2^2 \quad 1 \leq j \leq m$$

where  $A$  is orthonormal basis for  $\text{range}(\mathcal{A})$

- Rotated subspace:  $\text{range}(\mathcal{B})$ ,  $\mathcal{B}$  is  $m \times n$  with  $\text{rank}(\mathcal{B}) = n$   
Perturbed leverage scores

$$\ell_j(\mathcal{B}) = \|e_j^T \mathcal{B}\|_2^2 \quad 1 \leq j \leq m$$

where  $B$  is orthonormal basis for  $\text{range}(\mathcal{B})$

Question: How close are  $\ell_j(\mathcal{B})$  to  $\ell_j(\mathcal{A})$ ?

# Principal Angles between Column Spaces

$A$  and  $B$  are  $m \times n$  with orthonormal columns

- SVD of  $n \times n$  matrix  $A^T B = U \Sigma V^T$

$$\Sigma = \text{diag}(\cos \theta_1 \quad \cdots \quad \cos \theta_n)$$

- Principal angles  $\theta_j$  between  $\text{range}(A)$  and  $\text{range}(B)$

$$1 \geq \cos \theta_1 \geq \dots \geq \cos \theta_n \geq 0$$

$$0 \leq \theta_1 \leq \dots \leq \theta_n \leq \pi/2$$

- Special cases

If  $\text{range}(A) = \text{range}(B)$  then  $\Sigma = I_n$  and all  $\theta_j = 0$

If  $A^T B = 0$  then  $\Sigma = 0$  and all  $\theta_j = \pi/2$

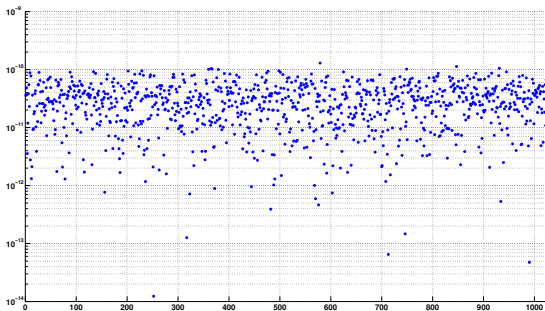


# Uniform Leverage Scores

$A$  is  $m \times n$  Hadamard  $m = 1024, n = 50$ , all leverage scores  $\ell_j(A) = n/m$

Angles:  $\cos \theta_1 = 1$   $\sin \theta_n \approx 10^{-8}$

Absolute errors  $|\ell_j(B) - \ell_j(A)|$  vs index  $j$



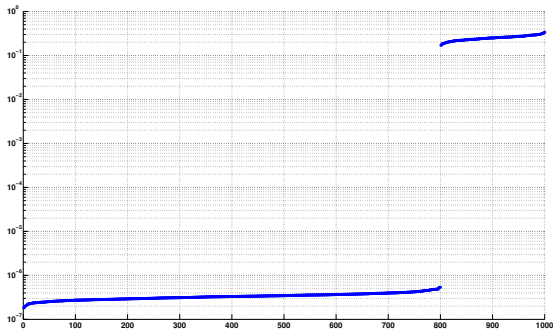
Absolute errors  $\leq 10^{-10}$

# 20% Large Leverage Scores

$A$  is  $m \times n$   $m = 1000, n = 50$

800 small  $\ell_j(A) \leq 10^{-6}$     200 large  $\ell_j(A) \approx .3$

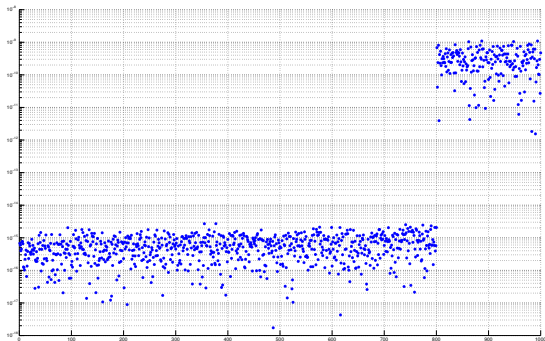
Leverage scores  $\ell_j(A)$  vs index  $j$



# 20% Large Leverage Scores: Absolute Errors

Angles:  $\cos \theta_1 = 1$   $\sin \theta_n \approx 10^{-8}$

Absolute errors  $|\ell_j(B) - \ell_j(A)|$  vs index  $j$



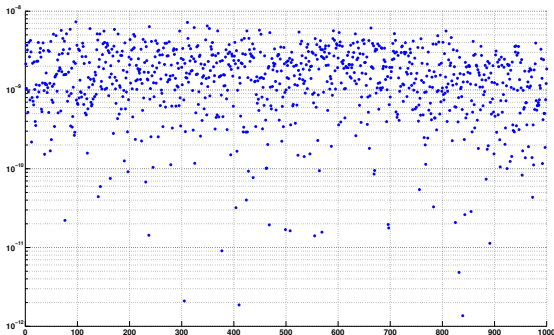
Small leverage scores: absolute errors  $\leq 10^{-15}$

Large leverage scores: absolute errors  $\leq 10^{-9}$

# 20% Large Leverage Scores: Relative Errors

Angles:  $\cos \theta_1 = 1$   $\sin \theta_n \approx 10^{-8}$

Relative errors  $|\ell_j(B) - \ell_j(A)|/|\ell_j(A)|$  vs index  $j$



All leverage scores have relative errors  $\leq 10^{-8}$

# Sensitivity of Leverage Scores to Subspace Rotation

$A$  and  $B$  are  $m \times n$  with orthonormal columns

- Angles between  $\text{range}(A)$  and  $\text{range}(B)$

$$0 \leq \theta_1 \leq \dots \leq \theta_n \leq \pi/2$$

- Perturbation bounds

$$l_j(B) \leq \left( \cos \theta_1 \sqrt{l_j(A)} + \sin \theta_n \sqrt{1 - l_j(A)} \right)^2$$

$$l_j(A) \leq \left( \cos \theta_1 \sqrt{l_j(B)} + \sin \theta_n \sqrt{1 - l_j(B)} \right)^2 \quad 1 \leq j \leq m$$

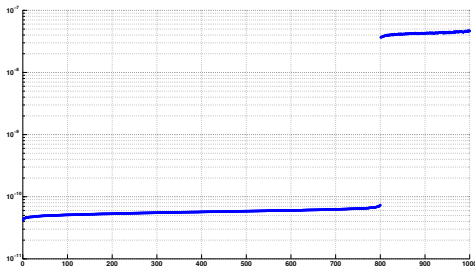
Leverage scores of  $A$  and  $B$  are **close**,  
if **all angles** between  $\text{range}(A)$  and  $\text{range}(B)$  are **small**

# Tightness of Bound

$A$  is  $m \times n$   $m = 1000, n = 50$   $\sin \theta \approx 10^{-8}$

800 small  $\ell_j(A) \leq 10^{-6}$     200 large  $\ell_j(A) \approx .3$

Bound -  $\ell_j(B)$  vs index  $j$



Bound tight in relative sense:  $\text{Bound} - \ell_j(B) \lesssim 10^{-8} \ell_j(B)$

## Large Leverage Scores

$A$  and  $B$  are  $m \times n$  with orthonormal columns

- Perturbed leverage score is large:  $l_k(B) \geq 1/2$  for some  $k$

$$\begin{aligned} \frac{l_k(A)}{(\cos \theta_1 + \sin \theta_n)^2} &\leq l_k(B) \\ &\leq \left( \cos \theta_1 \sqrt{l_k(A)} + \sin \theta_n \sqrt{1 - l_k(A)} \right)^2 \end{aligned}$$

- Exact leverage score is also large:  $l_k(A) \geq 1/2$

$$\frac{l_k(A)}{(\cos \theta_1 + \sin \theta_n)^2} \leq l_k(B) \leq (\cos \theta_1 + \sin \theta_n)^2 l_k(A)$$

Upper and lower bounds for large leverage scores

# Summary: Sensitivity to Subspace Rotation

Leverage scores of matrices with orthonormal columns

- Leverage scores are **close**, if subspace rotation **small**
- **Small** leverage scores as **insensitive** as **large** ones
- Our perturbation bounds are qualitatively **informative**
- **Simpler** bounds for **special cases**:  
Large leverage scores,  $m = 2n$



## Sensitivity of Leverage Scores to Matrix Perturbations

# Norm-wise Relative Matrix Perturbations

- $\mathcal{A}$  and  $\mathcal{A} + \mathcal{E}$  are  $m \times n$  of rank  $n$
- Two-norm condition number and relative perturbation

$$\kappa \equiv \|\mathcal{A}\|_2 \|\mathcal{A}^\dagger\|_2 \quad \epsilon \equiv \|\mathcal{E}\|_2 / \|\mathcal{A}\|_2$$

- First-order bounds

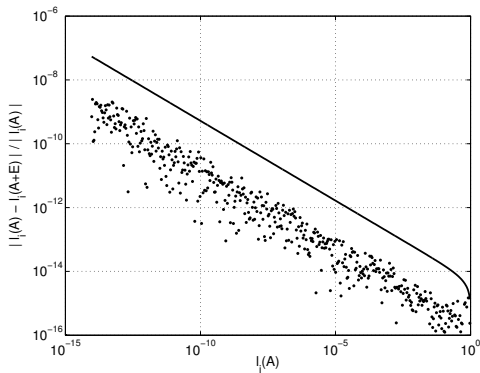
$$\left| \frac{\ell_j(\mathcal{A} + \mathcal{E}) - \ell_j(\mathcal{A})}{\ell_j(\mathcal{A})} \right| \leq 2 \sqrt{\frac{1 - \ell_j(\mathcal{A})}{\ell_j(\mathcal{A})}} \kappa \epsilon + \mathcal{O}(\epsilon^2)$$

- Sensitivity proportional to condition number of  $\mathcal{A}$   
Leverage scores sensitive if  $\kappa \gg 1$
- Small leverage scores more sensitive than large ones

# Small Perturbations

$\mathcal{A}$  is  $m \times n$   $m = 500$ ,  $n = 15$ ,  $\kappa = 1$ ,  $\epsilon \approx 10^{-15}$

Relative error  $|\ell_j(\mathcal{A} + \mathcal{E}) - \ell_j(\mathcal{A})| / \ell_j(\mathcal{A})$  vs  $\ell_j(\mathcal{A})$

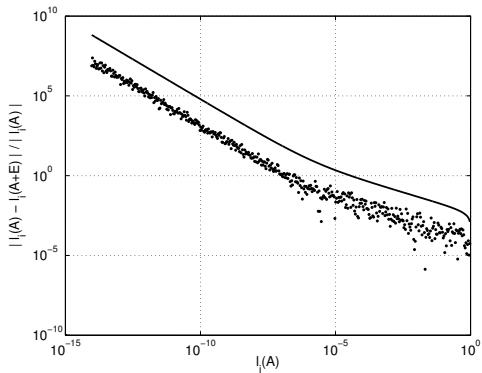


Bound reflects behavior of error

# Large Perturbations

$\mathcal{A}$  is  $m \times n$   $m = 500$ ,  $n = 15$ ,  $\kappa = 1$ ,  $\epsilon \approx 10^{-3}$

Relative error  $|\ell_j(\mathcal{A} + \mathcal{E}) - \ell_j(\mathcal{A})| / \ell_j(\mathcal{A})$  vs  $\ell_j(\mathcal{A})$



Bound reflects behavior of error

# Summary

## Motivation

Sampling *rows* from matrices  $Q$  with *orthonormal columns*

Want: *Condition number of sampled matrix*  $\kappa(SQ)$

Condition number depends on *leverage scores*

Largest leverage scores = *coherence*

## Matlab App `kappaSQ` (arXiv:1402.0642)

Sampling & bounds in *non-asymptotic context*

Compare different *probabilistic bounds*

Run experiments and test *tightness* of bounds

*"Publication-ready" plots*

## Sensitivity of Leverage Scores, to:

- Subspace rotations

Leverage scores *close* if subspace rotation *small*

Small leverage scores *as insensitive* as large ones

- Relative matrix perturbations

Sensitivity depends on *condition number* of matrix

Small leverage scores *more sensitive* than large ones